

# Ultra-Fast All-Optical Logic Circuit With The Massive Use of Non-Linear Material

Samir Sahu<sup>1\*</sup> and Shantanu Dhar<sup>2</sup>

<sup>1</sup>Department of Physics and Technophysics, Vidyasagar University, Midnapore, INDIA

\*Corresponding Author E-mail: [tosamirsahu@gmail.com](mailto:tosamirsahu@gmail.com)

<sup>2</sup>Department of Physics, Jhargram Raj College, Jhargram, INDIA

**Abstract:** Nonlinear material based all-optical switching mechanism is used here to develop the all-optical logic operation scheme. A single bit logic unit has been accessed first and is elevated to a higher bit logic unit in course. These circuits can execute innumerable logic operations and remarkably, they are all-optical and fully parallel in nature. These all-optical logic units can gear up to the highest capability of optical performance in high-speed all-optical computers.

**Keywords:** Nonlinear material, optical switching, all-optical NOT gate, multi-output gate, multiplexing circuit, logic circuit

## I. INTRODUCTION

No day can be imagined today without the talismanic touch of computer. In the coming future, we are to face rising usage of the Internet along with more complex and sophisticated problems in different walks of life, as computer is widening its globe. It compels us to think of such a computer which is not only faster but also more reliable. Unfortunately, the VLSI technology, so far is being used in electronics, is already in a saturated state from both the aspects – size and speed. To deal with the soaring demands, optics has already established its validity in various arithmetic [1-6], logic [7-13], algebraic [14-18], and image operations [19-22] in last few years, preliminary because of its promising features of parallel, high speed, high bandwidth non-interfering communications [5, 7, 9-11, 19, 23-24]. These gifted advantages lead us to our dream goal of making the fastest possible computer, specifically, a super fast optical computer that will outperform the fastest possible electronic computer.

Different techniques have been proposed and developed to implement several components of super fast computer in optical domain. In the recent past, all-optical switching mechanism [5, 9-13, 25-27, 29, 34] by nonlinear optical material is well-known as one of such promising techniques. Arithmetic operations as well as logical operation are the two essential tasks in any type of computing system. The Arithmetic operation scheme [6] has been proposed earlier. This paper proposes a scheme for the all-optical implementation of logic operational circuit with proper use of nonlinear material-based all-optical switching mechanism. At first, we design an all-optical gate (Multi-output gate) which can give three distinct logic outputs of the very same inputs at the same time. Then, an all-optical multiplexure is designed. Finally, we combine these two circuits to develop a single bit logic operation scheme. As the circuit is purely all-optical in nature, it is very simple and very fast. All the basic logic operations and XOR operation can be achieved with extreme accuracy. The scheme can be extended to a higher bit logic operation scheme easily. An ALU of the proposed goal, an optical computer, can be implemented through this scheme.

## II. ULTRA-HIGH-SPEED ALL-OPTICAL SWITCHING BEHAVIOR OF NONLINEAR MATERIAL AND ITS USES AS ALL-OPTICAL NOT GATE, MULTI-OUTPUT GATE AND AND GATE

The phenomenon photorefractivity [10-12, 28] of some nonlinear optical material is used in nonlinear all-optical intensity switching mechanism. The refractive index [10-12, 20, 28, 30-31] of some nonlinear materials (NLM) such as carbon disulfide, pure silica, potassium dihydrophosphate (KH<sub>2</sub>PO<sub>4</sub> (KDP) crystal etc. varies linearly with the intensity of the light incident on it. The refractive index ( $n$ ) of such isotropic dielectric non-crystalline media can be put into an equation as (1). Here  $n_0$  is the linear term,  $n_1$  is the nonlinear correction term and  $I$  is the intensity of the incident light beam on the material.

$$n = n_0 + n_1 I$$

### A. Ultra Fast All-Optical Switching Behavior of Nonlinear Material

We can implement the switching mechanism with such nonlinear material by taking an interface between two media of which one is a linear material (LM), whose refractive index  $n_0$  is independent of the intensity of light and the other is aforesaid NLM. A laser beam, highly intense polarized light, preferably pulse laser of intensity  $I_1$ , is allowed to incident on the interface from linear to nonlinear part in a particular direction XO (incidence angle  $\theta_1$ ) as it depicted in Fig. 1. The refracted beam from the NLM follows the path OZ. But when another higher intense laser beam of intensity  $I_2$  ( $I_2 > I_1$ ) is made to incident along XO, after refraction from the NLM the light passes through OY direction (angle of refraction  $\theta_2$ ). The deviation of refractive angle for different incident light intensity  $I_1$  and  $I_2$  is  $\angle ZOY = \Delta\theta_2$ . Thus the combination of LM and NLM may act nicely as a directional all-optical switch. This is the unit block of our proposed logic circuit.

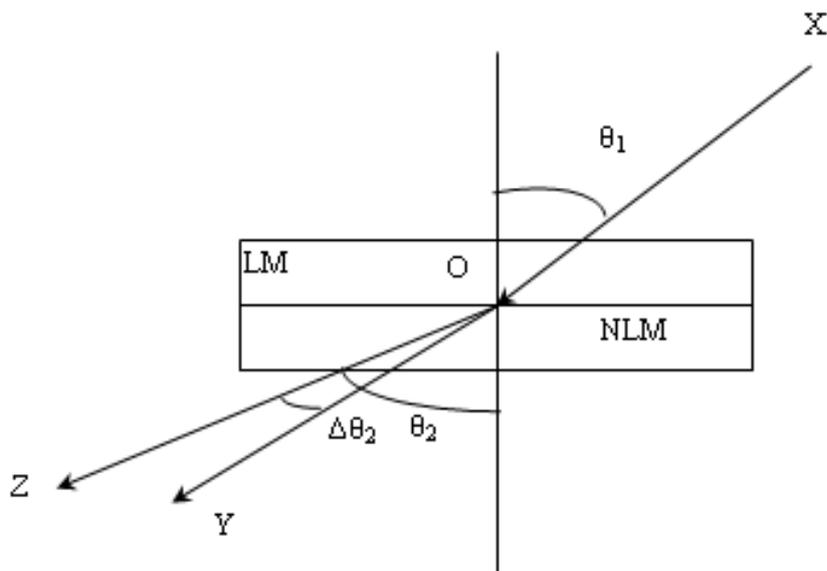


Figure 1: Intensity switching of optical nonlinear material

In the expression of refractive index in Eq. (1)  $n_0$  is linear term and  $n_1$  is the nonlinear correction term. For carbon disulfide [10-12, 29] (CS<sub>2</sub>)  $n_0 = 1.63$ ,  $n_1 = 514 \times 10^{-20} \text{ m}^2/\text{W}$ . and for fused silicon dioxide [10-12, 29] (SiO<sub>2</sub>)  $n_0 = 1.458$ ,  $n_1 = 2.7 \times 10^{-20} \text{ m}^2/\text{W}$ . If we use CS<sub>2</sub> and SiO<sub>2</sub> as

nonlinear materials and the pulse laser of intensity  $I = 2 \times 10^{18} \text{ W/m}^2$  as a source, we can estimate the deviations of light in two cases as given in Table I.

Table I. Estimation of the deviation of pulsed laser light when passing through carbon disulfide (CS<sub>2</sub>) and silicon dioxide (SiO<sub>2</sub>)

| Material                             | Angle of incidence ( $\theta_1$ ) | Incident light intensity           | $n$<br>( $= n_0 + n_1 I$ ) | Angle of refraction ( $\theta_2$ ) | Deviation<br>( $\Delta\theta_2 = \theta'_2 - \theta''_2$ ) |
|--------------------------------------|-----------------------------------|------------------------------------|----------------------------|------------------------------------|--|
| Carbon disulfide (CS <sub>2</sub> )  | 45 deg                            | $I=2 \times 10^{18} \text{ W/m}^2$ | 11.91                      | 3.404 deg = $\theta'_2$            | 1.578 deg  |
|                                      | 45 deg                            | 2I                                 | 22.19                      | 1.827 deg = $\theta''_2$           |  |
| silicon di-oxide (SiO <sub>2</sub> ) | 45 deg                            | $I=2 \times 10^{18} \text{ W/m}^2$ | 1.512                      | 27.883 deg = $\theta'_2$           | 1.041 deg  |
|                                      | 45 deg                            | 2I                                 | 1.566                      | 27.842 deg = $\theta''_2$          |  |

The logic gates [10-12, 28-29] are implemented in optics using NLM by taking the presence of light signal as 1 and the absence of it as 0.

The implementation of such logic gates can be done by using some femtosecond laser pulses and 1-mm-thick KDP crystal at the peak intensity of  $0.6 \text{ TW/cm}^2$  and duration of 60 fs [10-12, 29]. M. Choi et al. show that a single-layer terahertz metamaterial [31] has a peak refractive index of 38.6 while maintaining low losses. It is a broadband, extremely high index of refraction going beyond the limit that is attainable with naturally existing substances, lead sulphide, and strontium titanate [31].

### B. All Optical NOT Gate

To implement an all-optical NOT gate using non-linear material, a constant intensity pulse laser source (CILS) is used as shown in Fig. 2(a). It is also called probe beam. Here  $A_1$  is taken as input beam. A detector is placed at  $D_3$  will detect the output beam after refraction. If  $A_1$  is absent, the light will follow a path  $OG_3$  and will be detected by the detector due to presence of CILS. But if  $A_1$  is present, after refraction, the light will follow a path other than  $OG_3$ , may be  $OG_4$ , and the detector will not detect any light signal. So  $D_3$  is equals to  $\overline{A_1}$ . Thus the system (NG block) acts as all-optical NOT gate.

### C. All Optical Multi-Output Gate

We design an all-optical Multi-output gate which can give three different logic outputs of same inputs at the same time, shown in Fig 2(b). Here  $A_1$  and  $B_1$  are two input channels. Three detectors placed at  $D_0$ ,  $D_1$  and  $D_2$  give the outputs. If any one of the input signals carries light, after refraction light will follow the path  $OG_1$ . The path  $OG_2$  will carry light if both  $A_1$  and  $B_1$  are present.

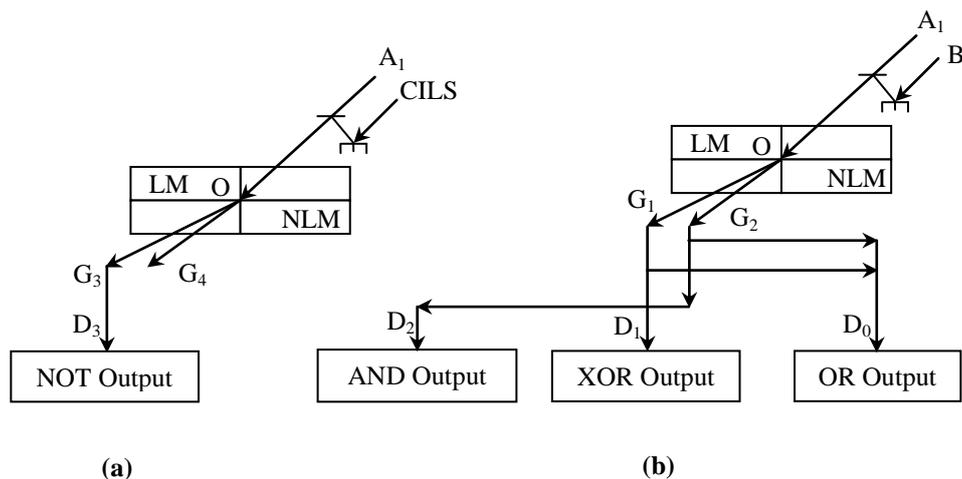


Figure 2: All-optical logic gates (a) NOT gate (b) Multi-Output gate

Now, we direct light from both the path  $OG_1$  and  $OG_2$  to  $D_0$  detector. When both the channels remain dark,  $D_0$  gives 0, otherwise it is 1. So, at  $D_0$ , we get simply logical OR Output of  $A_1$  and  $A_2$  inputs, i.e.  $D_0 = A_1 \vee A_2$ . If only one input carries light, the light ray after refraction will be detected by the detector at  $D_1$ , otherwise not. Thus,  $D_1$  gives  $(A_1 \oplus A_2)$  output. Now when both the channels carry light signal, the light beam after passing through the block will be detected by the detector at  $D_2$ , otherwise not. Hence one can get AND output of the inputs  $A_1$  and  $A_2$  at  $D_2$  terminal ( $D_2 = A_1 \wedge A_2$ ). So the block MG simultaneously acts as OR, XOR and AND gates (We would like to define it as Multi-output (MG) gate).

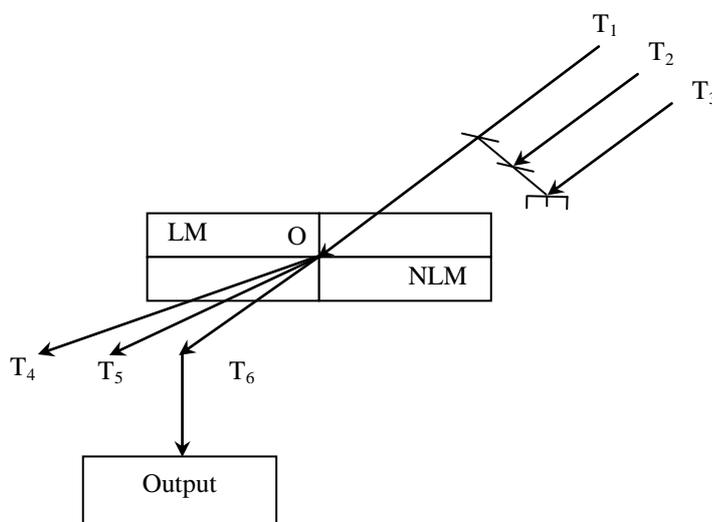


Figure 3: All-optical three input AND gate

D. All Optical AND Gate

The Three-input-all-optical AND gate using NLM is shown in Fig 3. Here  $T_1$ ,  $T_2$  and  $T_3$  are the three input channels. Light can travel through any of the three paths  $OT_4$ ,  $OT_5$ ,  $OT_6$ . A detector placed at  $T_6$  gives the output. Now when all the channels carry light signal, the light beam after refraction will be detected by the detector at  $T_6$ , otherwise not.

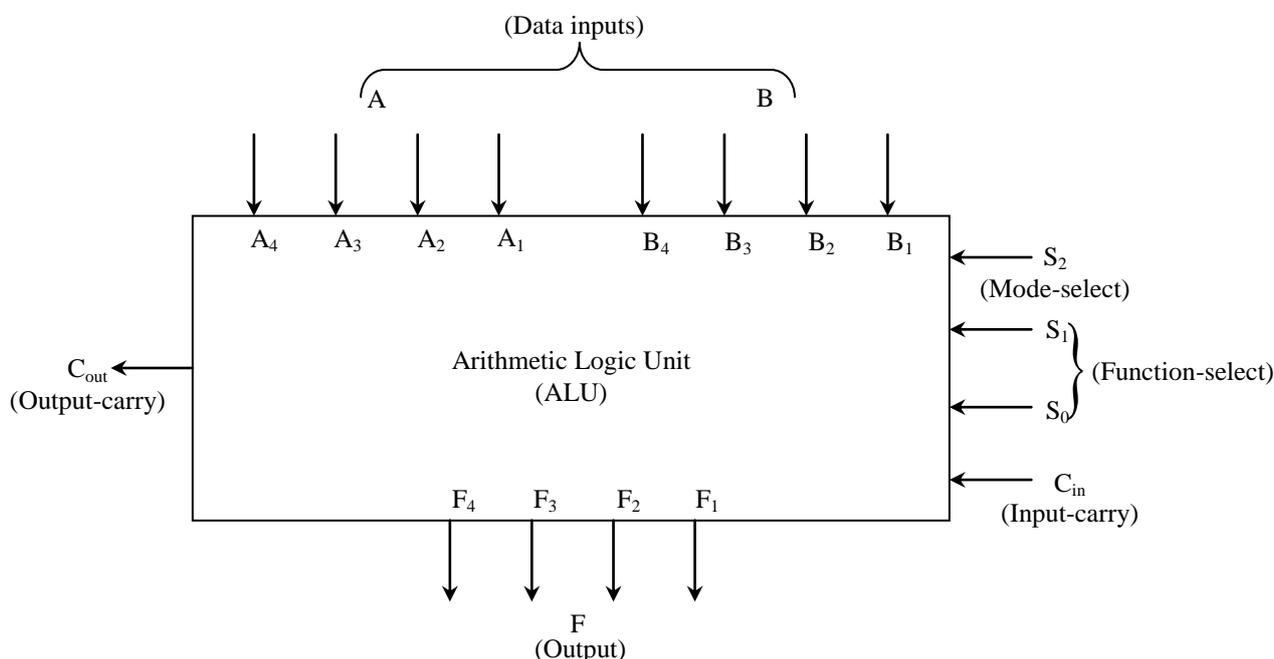


Figure 4: Block diagram of a conventional four-bit ALU

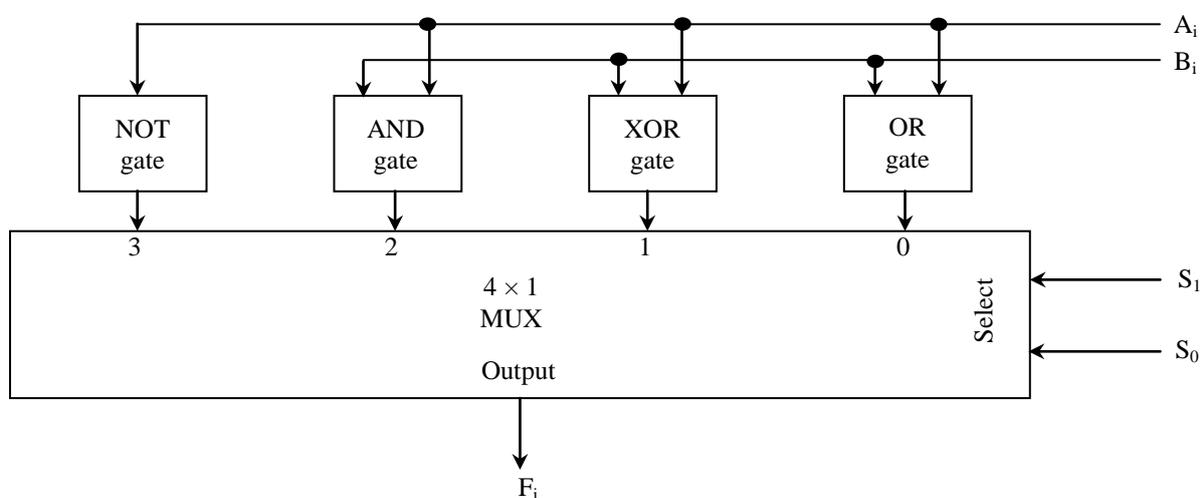


Figure 5: Electronically addressed one bit ( $i^{th}$  stage) logic circuit

### III. CONVENTIONAL ELECTRONIC LOGIC CIRCUIT IN ALU

The block diagram of a 4-bit conventional ALU [32-33] is shown in Fig. 4. If  $S_2 = 1$  ALU acts as logic unit, Input-Carry ( $C_{in}$ ) and Output-Carry ( $C_{out}$ ) have no meaning. The simplest conventional electronic logic circuit is shown in Fig. 5 which depicts one typical stage designated by subscript  $i$ . It can be repeated  $n$  times for an  $n$ -bit logic circuit. The output of four logic gates, OR, XOR, AND and NOT respectively, are the data inputs of a  $4 \times 1$  MUX. The output  $F_i$  is one of the four logic operations according to the proper combination of function-select inputs,  $S_1$  and  $S_0$ . The function table is shown in Table II.

Table II. Function table for the all-optical logic circuit

| <i>Function-selection variables</i> |       | <i>Output equals to</i> | <i>Circuit action</i> |
|-------------------------------------|-------|-------------------------|-----------------------|
| $S_1$                               | $S_0$ |                         |                       |
| 0                                   | 0     | $F_i = A_i \vee B_i$    | OR Operation          |
| 0                                   | 1     | $F_i = A_i \oplus B_i$  | XOR Operation         |
| 1                                   | 0     | $F_i = A_i \wedge B_i$  | AND Operation         |
| 1                                   | 1     | $F_i = \overline{A_i}$  | NOT Operation         |

### IV. CONCEPT OF MULTIPLEXING

Multiplexing is nothing but transmitting a large number of information units over a lesser number of outlines [5, 32-33]. It is also named as data selector, because it selects one of many inputs and guides the binary information to the output channel. The digital multiplexer is a combinational circuit that selects binary information from one of several input lines and directs it to a single output channel. The choice of a particular input line is restricted by a set of selection input lines. Normally, there is  $2^n$  input channels and the bit combinations of  $n$  selection inputs determine that very input which is selected. The block diagram of a 4-line to 1-line multiplexer ( $4 \times 1$  MUX) is shown in Fig. 6(a). Each of the four data input lines,  $D_0$  to  $D_3$ , is selected and sent to the output  $Y$  by the proper combination of select inputs  $S_1$  and  $S_0$ . The function table is shown in Table III.

Table III. Function table for of  $4 \times 1$  multiplexure circuit

| <i>select inputs</i> |       | <i>Output Y equals to</i> |
|----------------------|-------|---------------------------|
| $S_1$                | $S_0$ |                           |
| 0                    | 0     | $D_0$                     |
| 0                    | 1     | $D_1$                     |
| 1                    | 0     | $D_2$                     |
| 1                    | 1     | $D_3$                     |

## V. ALL-OPTICAL MULTIPLEXING SYSTEM

The scheme of all-optical multiplexer [25, 32-33] is made of with the combinational blocks of linear and nonlinear materials shown in Fig. 6(b). Here NG1 and NG2 are two all-optical NOT gates and AG1 to AG4 are four three input AND gates. The input lines  $S_1$  and  $S_0$  act as select inputs of the data selector.  $D_0$ ,  $D_1$ ,  $D_2$  and  $D_3$  are the four data input lines and  $Y$  is the final output of the MUX.

Here  $S_0$ , is the input of NG1 and  $S_1$  is the input of NG2. CILS1 and CILS2 are the two probe beams of the two NOT gates NG1 and NG2 respectively. As NG1 and NG2 are two NOT gates,  $S_0$  and  $S_1$  are inverted through  $O_1C_1$  and  $O_2C_2$  paths respectively and they are represented as  $S'_0$  and  $S'_1$ . We take the direct input  $S_0$  and also the complement of  $S_0$  ( $S'_0$ ) as two outputs from NG1. Similarly, from NG2,  $S_1$  and  $S'_1$  are in use.

Now, AG1, the first AND gate, has inputs  $D_0$ ,  $S'_0$  and  $S'_1$  and here we pick up the output  $N_3$  solely. The first data input,  $D_0$  is transferred to the output  $Y$  through AG1 if and only if  $S'_0$  and  $S'_1$  are both at high state (i.e.  $S_0S_1 = 00$ ). In this situation exclusively the block AG1 will be active but not the other blocks AG2, AG3 and AG4. Secondly  $D_1$ ,  $S_0$  and  $S'_1$  are the inputs of AG2 block and here the output is taken from  $N_4$  terminal. At the third AND gate, AG3 gets inputs from  $D_2$ ,  $S'_0$  and  $S_1$  and the output is received from  $N_5$  end.  $D_3$ ,  $S_0$  and  $S_1$  are taken as the three inputs of AG4 and here  $N_6$  channel yields the output. Similar to the first case at AG1, the data at  $D_1$ ,  $D_2$  and  $D_3$  are transferred to the output  $Y$  respectively if and only if other two inputs of each AND gates are at logical '1' stage. Ultimately, all the outputs ( $N_3$ ,  $N_4$ ,  $N_5$  and  $N_6$ ) of four AND gates are united to form the ultimate output  $Y$  of the  $4 \times 1$  MUX.

Now, we need to explain the operation of the multiplexer referred to Fig. 6(b). Let us take the select input  $S_0 = S_1 = 0$  (i.e.  $S_0$  and  $S_1$  both remain dark). Due to presence of probe beams CILS1 and CILS2, light signal will follow the path  $O_1C_1$  and  $O_2C_2$  when passing through the block NG1 and NG2 respectively. At present, we have  $S_0 = S_1 = 0$  and  $S'_0 = S'_1 = 1$ . In this condition, the first AND gate (AG1) is active to transfer output through  $O_3N_3$  to  $Y$ . We can get light at  $Y$  ( $=1$ ) when  $D_0 = 1$ , and we will not have it when  $D_0 = 0$ . In the present situation, all other AND gates will stay inactive, because either one or both of the control inputs  $S_0$  and  $S_1$  have no light. So,  $N_4$ ,  $N_5$  and  $N_6$  terminals have no light signal whatever may be the status of  $D_1$ ,  $D_2$  and  $D_3$  lines. Thus, we come to the decision that only the data input  $D_0$  is connected to the output channel  $Y$  in absence of  $S_0$  and  $S_1$ . Similarly, if  $S_0 = 1$  and  $S_1 = 0$ , only the AG2 gate will be active and none other than  $D_1$  is connected to  $Y$ . In the same manner, when  $S_0 = 0$  and  $S_1 = 1$ , only the  $D_2$  channel will remain connected to  $Y$ , and when  $S_0 = S_1 = 1$ , only  $D_3$  will be connected to the final output.

## VI. ALL-OPTICAL LOGIC CIRCUIT

We now propose a simplest all-optical logic operation scheme which is illustrated in Fig. 7. This picture expresses one typical stage chosen by subscript  $i$ . It can be repeated  $n$  times to obtain an  $n$ -bit logic circuit. The output of four logic gates, OR, XOR, AND and NOT respectively, are the data inputs of a  $4 \times 1$  MUX. According to the proper combination of function-select inputs,  $S_1$  and  $S_0$ , the output  $F_i$  is one of the four logic operations. The function table is shown in Table II.

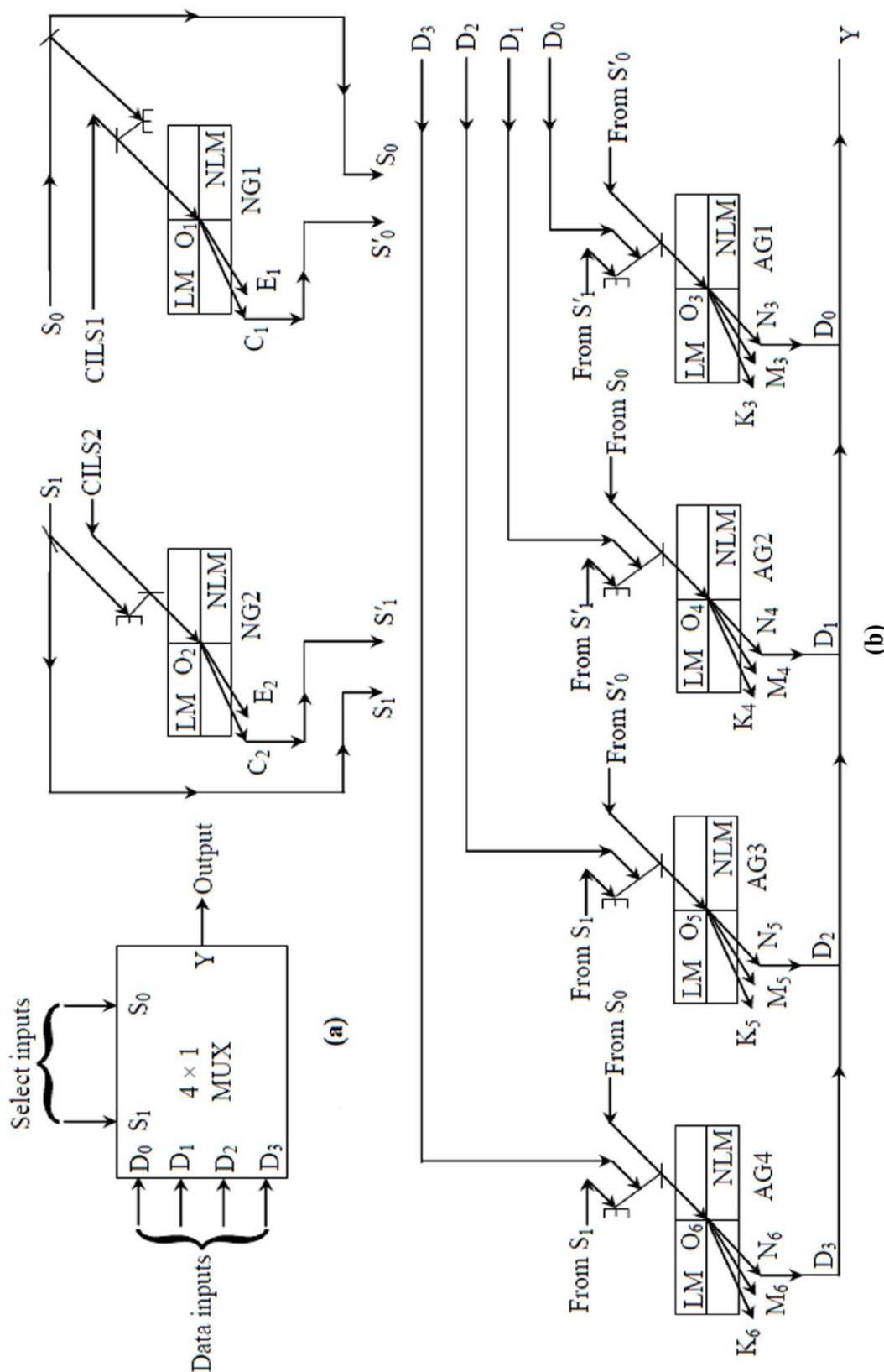


Figure 5: Multiplexer (a) Block diagram of a conventional 4-line to 1-line multiplexer (b) All-optical multiplexer circuit.

Here  $A_i$  and  $B_i$  are the two one bit ( $i^{\text{th}}$  bit of an  $n$  bit binary numbers) binary numbers used as data inputs of our scheme which are combined with each other to generate a logic operation at  $F_i$  output. The two selection variables,  $S_1$  and  $S_0$ , specify the particular logic operation to be produced at the final output channel  $F_i$ . It is possible to create a total four logic operations (OR, XOR, AND and NOT) by the combination of two control bits  $S_1$  and  $S_0$ . The circuit is constructed by eight blocks (combinations of linear and nonlinear materials), our unit blocks. Among eight unit blocks, NG1, NG2 and NG3 are the three NOT gates, AG1, AG2, AG3 and AG4 are the four three input AND gates and MG is a Multi-output gate.

To explain it more legibly, let us divide the circuit into two parts.

**Part I: Logic Operation Part:** The Logic Operation Part is enclosed by the dotted rectangle in Fig. 7. The blocks, MG and NG3, form the first part.  $A_i$  and  $B_i$  are inputs to MG, the Multi-output gate. The three outputs from it are  $D_0 (=A_i \vee B_i)$ ,  $D_1 (=A_i \oplus B_i)$ ,  $D_2 (=A_i \wedge B_i)$ . The NOT gate has input  $A_i$  along with the probe beam CILS. One can expect the output of NG3,  $D_3 = \overline{A_i}$ .

**Part II: 4×1 Multiplexer Part:** The rest part of Fig. 7 other than the Logic Operation Part is 4×1 Multiplexer Part. Two NOT gates, NG1 and NG2, and four AND gates, AG1, AG2, AG3 and AG4, construct the multiplexer part of the logic circuit. The select variables  $S_1$  and  $S_0$  are inverted by NG2 and NG1 respectively. The combination of two variables ( $S'_0, S'_1$  for AG1;  $S_0, S'_1$  for AG2;  $S'_0, S_1$  for AG3 and  $S_0, S_1$  for AG4) is used in the said manner as two inputs of the four three inputs AND gates. The outputs from Logic Operation Part,  $D_0, D_1, D_2$  and  $D_3$ , are third inputs of each AND gates AG1 to AG4 respectively.

Now four cases may arise

**Case 1:** When  $S_1 = S_0 = 0$ , then  $S'_1 = S'_0 = 1$ , that activates none other than the first AND gate AG1. As a result, only  $D_0 (=A_i \vee B_i)$  is transferred to the final output  $F_i$  through  $O_3N_3$  and  $Y_0$ . In this situation, other three AND gates (AG2, AG3 and AG4) remain inactive, independent to the input variables. Thus we can conclude that the logic operation scheme produces the final output  $F_i = A_i \vee B_i$  when  $S_0$  and  $S_1$  both at logical '0' state. In other words, we can say that the circuit performs the OR logic operation for  $S_1S_0 = 00$ .

**Case 2:** Now when  $S_0 (= 1)$  carries light but  $S_1 (= 0)$  does not, we have  $S'_1 = 1$  and  $S'_0 = 0$ . This turns the second AND gate, AG2 active but not the other three AND gates. As an outcome  $F_i$  is nothing but  $D_1 (=A_i \oplus B_i)$ . The circuit gives EX-OR operation when  $S_1 = 0, S_0 = 1$ .

**Case 3:** Next, when  $S_0$  is dark but  $S_1$  is lighted, i.e.  $S_1 = 1$  and  $S_0 = 0$ , then  $S'_1 = 0$  and  $S'_0 = 1$ . As  $S_1$  and  $S'_0$  both carry light signal, AG3 becomes active but not AG4, AG2 and AG1. Only the  $D_2 (=A_i \wedge B_i)$  terminal is supposed to be connected to  $F_i$  via  $O_5N_5$  and  $Y_2$ . The ultimate result of our scheme,  $F_i$  generates AND operation of the data inputs  $A_i$  and  $B_i$  if and only if  $S_1$  is equal to 1 and  $S_0$  is equal to 0.

**Case 4:** In the penultimate condition, when both the select bits  $S_0$  and  $S_1$  are present ( $S_1 = S_0 = 1$ ), there will be on light at  $S'_1$  and  $S'_0$  ends. Now, the fourth AND gate (AG4)

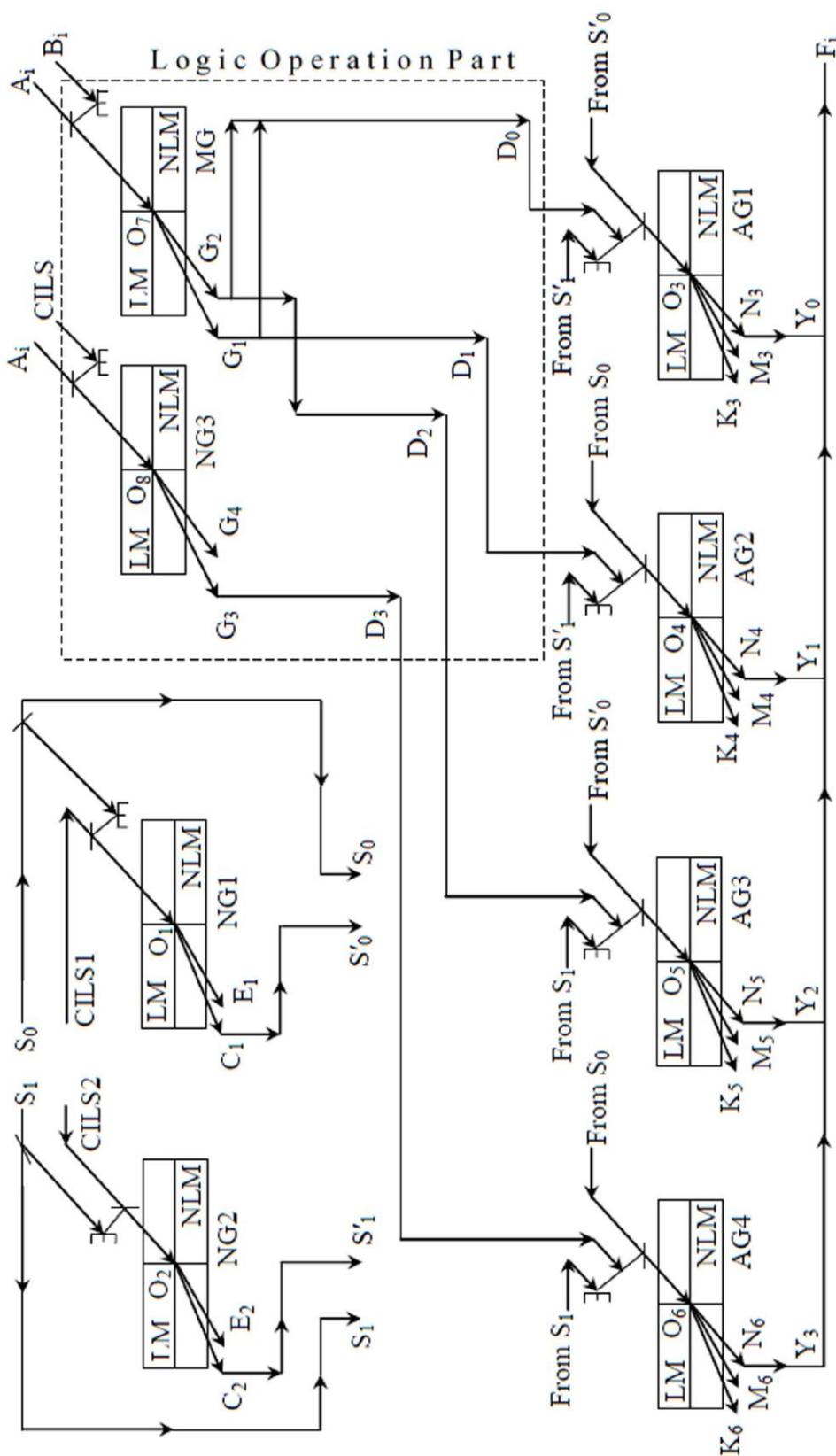


Figure 7: All-optical one bit ( $i^{th}$  stage) logic circuit.

shifts  $D_3 (= \overline{A_i})$  output to  $F_i$  through  $O_6N_6$ . All the other AND gates are inactive now, since either one or both of the control inputs  $S'_0$  and  $S'_1$  have no light. Thus we come to the decision that in presence of  $S_0$  and  $S_1$  (i.e.  $S_1S_0 = 11$ ) only the NOT operation of  $A_i$  is executed by the logic circuit.

## VI. CONCLUSION

The proposed technique of all-optical implementation of logic operation scheme is very fast (above THz) as it is fully all-optical. The light signals that are severally used, bended and the feedback light signals from the outputs are made by mirrors and beam splitters to make the circuits simple. This operation scheme should be the second step on our dream way to all-optical Arithmetic and Logic Unit. Along with this the circuit being parallel becomes remarkably fast. Proper findings of nonlinear material may be a significant issue here. Essentially inputs should be chosen properly for proper function of the system.

## VII. REFERENCES

- [1] S. Mukhopadhyay, A. K. Datta and A. Basuray (1990). A new technique of arithmetic operation using positional residue system. *Appl. Opt.* vol. 29(20), pp. 2981-2983.
- [2] A. Gonzalez-Marcos and J. A. Martin-Pereda (2001). Method to analyze the influence of hysteresis in optical arithmetic units. *Opt. Eng.* vol. 40, pp. 2371-2385.
- [3] S. Sahu, R. R. Pal and S. Dhar (2011). Nonlinear Material Based All-Optical Parallel Subtraction Scheme: an Implementation. *International J. of Optoelectronics Engineering.* vol. 1(1), pp. 7-11.
- [4] R. S. Fyath, A. A. W. Alsaffar and M. S. Alam (2002). Optical binary logic gate-based modified signed-digit arithmetic. *Opt. Laser Technol.* vol. 34, pp. 501-508.
- [5] K. R. Choudhury and S. Mukhopadhyay (2004). Binary optical arithmetic operation scheme with tree architecture by proper accommodation of optical nonlinear materials. *Opt. Eng.* vol. 43(1), pp. 132-136.
- [6] S. Sahu and S. Dhar (2014). All-Optical Implementation of Arithmetic Operation Scheme using Optical Nonlinear Material Based Switching Technique. *Photonics and Optoelectronics (P&O).* vol. 3, pp. 37-50.
- [7] S. Randel, A. M. de Melo, K. Petermann, V. Marembert and C. Schubert (2004). Novel scheme for ultrafast all-optical XOR operation. *J. Lightwave Technol.* vol. 22, pp. 2808-2815.
- [8] X. Yang, C. Zhang, S. Qi, K. Chen, J. Tian and G. Zhang (2005). All-optical Boolean logic gate using azo-dye doped polymer film. *Optik* vol. 116, pp. 251-254.
- [9] S. Dhar and S. Mukhopadhyay (2005). All-optical implementation of ASCII by use of nonlinear material for optical encoding of necessary symbols. *Opt. Eng.* vol. 44(6), pp. 065201-1-5.
- [10] S. Dhar and S. Sahu (2008). All-optical implementation of S-R, clocked S-R and D flip-flops using nonlinear material. *Opt. Engineering.* vol. 47(6), pp. 065401-1-6.
- [11] S. Sahu and S. Dhar (2009). Implementation of clocked J-K, T and J-K Master Slave flip-flops with nonlinear material in All-optical Domain. *Opt. Engineering.* vol. 48(7), pp. 075401-1-7.
- [12] S. Sahu, R. R. Pal and S. Dhar (2011). A Novel Method of Implementing Nonlinear Material Based All-Optical Binary Half Subtractor and Full Subtractor System. *J. of Electron Devices.* vol. 10, pp. 493-498.
- [13] S. Sahu, R. R. Pal and S. Dhar (2011). Ultra-High Speed All-Optical T Flip-Flop Without Preset and Clear Using Non-Linear Material: a Theoretical Study. *J. of Phys. Sc.* vol. 15, pp. 241-250.
- [14] J. Y. Jau, F. Kiamilev, Y. Fainman, S.C. Esener and S. H. Lee (1988). Optical expert system based on matrix-algebraic formulation. *Appl. Opt.* vol. 27(24), pp. 5170-5175.
- [15] A. Andreoni, M. Bondani, M. A. C. Potenza, Y. N. Denisyuk and E. Puddu (2001). Boolean algebra operations performed on optical bits by the generation of holographic fields through second-order nonlinear interactions. *Rev. Sci. Instrum.* vol. 72, pp. 2525-2531.

- [16] S. Sahu, R. R. Pal and S. Dhar (2011). Implementation of 1-Bit random Access Memory Cell in All-Optical Domain with Non-linear material. *International J. of Optics and Application*. vol. 1(1), pp. 8–12.
- [17] S. Sahu, R. R. Pal and S. Dhar (2011). TeraHertz All-Optical Binary Register using D flip-flop with Non-linear Material: A Proposal. *J. of Electron Devices*. vol. 11, pp. 588–595.
- [18] S. Sahu, R. R. Pal and S. Dhar (2011). All-Optical Binary Counter by using T flip-flop: An Implementation. *International J. of Engineering, Science and Technology*. vol. 3(10), pp. 7799–7807.
- [19] Z. Chen and M. A. Karim (1999). Speed limitation of hybrid optical/digital sequential image processing. *Opt. Commun.* vol. 168(1-4), pp. 75–83.
- [20] N. Pahari and S. Mukhopadhyay (2005). An all optical R-S flip-flop by optical non-linear material. *J. of Opts.* vol. 34 (3), pp. 108–114.
- [21] A. Sinha and K. Singh (2003). A technique for image encryption using digital signature. *Opt. Commun.* vol. 218(4-6), pp. 229–234.
- [22] N. Nishimura, Y. Awatsuji and T. Kubota (2004). Two-dimensional arrangement of spatial patterns representing numerical data in input images for effective use of hardware resources in digital optical computing system based on optical array logic. *J. Parallel Distribut. Comput.* Vol. 64, pp. 1027–1040.
- [23] H. J. Caulfield, C. S. Vikram and A. Zavalin (2006). Optical logic redox. *Optik*. vol. 117, pp. 199–
- [24] A. K. Datta and S. Munshi (2007). Optical implementation of flip-flops using single-LCD panel. *Optical & Laser Technology*, vol. 39(9), pp. 2321-2329.
- [25] K. R. Choudhury, P. P. Das and S. Mukhopadhyay (2005). All-optical time-domain multiplexing-demultiplexing scheme with nonlinear material. *Opt. Eng.* vol. 44(3), pp. 035201-1–4.
- [26] S. Dhar and S. Mukhopadhyay (2006). All-optical decoding method for ASCII-coded data using nonlinear-material-based switching. *Opt. Eng.* vol. 45(11), pp. 115201-1–5.
- [27] K. R. Choudhury and S. Mukhopadhyay (2003). A new method of binary addition scheme with massive use of optical nonlinear material based system. *Chin. Opt. Lett.* vol. 1(3), pp. 040241-01–02.
- [28] D. Arivouli (2001). Fundamentals of optical nonlinear materials. *Pramana*. vol. 57(5,6), pp. 871– 883.
- [29] D. Samanta and S. Mukhopadhyay (2007). A method of maintaining the intensity level of a polarization encoded light signal. *J. of Phys. Sc.* vol. 11, pp. 87–91.
- [30] S. Mironov, V. Lozhkarev, V. Ginzburg and E. Khazanov (2009). High-efficiency second-harmonic generation of superintense ultrashort laser pulses. *Appl. Opt.* vol. 48, pp. 2051– 2057.
- [31] M. Choi, S. H. Lee, Y. Kim, S. B. Kang, J. Shin, M. H. Kwak, K.Y. Kang, Y. H. Lee, N. Park and B. Min (2011). A terahertz metamaterial with unnaturally high refractive index. *Nature*. vol. 470, pp. 369–373.
- [32] M. Morris Mano (2000). Digital Logic and Computer Design, 22<sup>nd</sup> ed., *Prentice-Hall of India Private Limited*, New Delhi, Chaps.5. 8 and 9.
- [33] R. P. Jain (2007). Modern Digital Electronics, 3<sup>rd</sup> ed., *Tata McGraw-Hill India*, New Delhi, Chaps. 6 and 13.
- [34] S. Mukhopadhyay (1990). An optical conversion system: from binary to decimal and decimal to binary. *Opt. Commun.* Vol. 76, pp. 309–312.